

A New Non-Orthogonal Space-Time Code with Low Decoding Complexity

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Abstract—A full diversity block space-time code over two transmit antennas and two symbol periods is introduced. In this method, each code is equal to the addition of two matrices; \mathbf{A}^m and $\mathbf{D}\mathbf{A}^n$, where \mathbf{A} and \mathbf{D} are the two constant matrices and m and n are the two data symbols. \mathbf{A} is selected such that the set $\mathbf{A}^m, 0 \leq m \leq 2^b - 1$ is closed under the matrix multiplication. This structure allows a simple Maximum Likelihood (ML) decoding method, and at the same time, simplifies the optimization of the coding advantage. Simulations show that the performance of the new code is very close to that of the Damen code [1] which is the best known block space-time code in terms of the coding advantage. Moreover, the decoding complexity of the proposed method is significantly lower than that of the Damen code.

Index Terms—Space-time coding, multiple-antenna arrays, transmit diversity, Maximum Likelihood (ML) decoding, wireless communications.

I. INTRODUCTION

RECENTLY, the use of multiple antennas at both ends of wireless links has received significant attention in order to increase the communication rate and/or to exploit the diversity [2]–[4]. Space-time code, a combination of modulation and channel coding across the space and the time dimensions, has been introduced by Tarokh *et al.* to take advantage of the spatial diversity [4]. Authors have introduced two criteria, based on the rank and the determinant of the substructure of the code matrices, to design space-time codes. These design criteria have been successfully used in other research [5]–[9]. Later, other design criteria have been developed for special scenarios such as low Signal to Noise Ratio (SNR), high data rates, and a large number of antennas [10]–[14].

Block orthogonal space-time coding is another method of designing full diversity space-time codes with an orthogonality property which simplifies decoding [5] [15]. However, it has been demonstrated that the use of an orthogonal structure sacrifices the achievable coding gain. Moreover, this structure is restrictive in the sense that it does not allow to achieve the channel capacity for more than one receive antenna [14]. A well known example of the orthogonal code is the so-called *Alamouti code* [15] which takes advantage of two transmit antennas over two symbol periods.

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Knowing the suboptimality of the Alamouti code and considering its practical significance, researchers have attempted to improve this code. For example, some have focused on improving the coding advantage of the orthogonal Space-Time Block Code (STBC) by concatenating the Alamouti scheme with a forward error correcting code [16]–[19]. Also, there have been some attempts to construct a code with a high rate by relaxing the orthogonality. In [20], a quasi-orthogonal code, based on using the Alamouti code as the building block is introduced. The rate of this code is one symbol Per Channel Use (PCU), however it offers a diversity which is only half of the maximum possible value. In [21], a full rate code is proposed to minimize a defined measure of the non-orthogonality. Non-orthogonality arises from increasing the rate above the maximum that is possible with an orthogonal structure.

There have been some attempts to increase the coding advantage by relaxing the orthogonality. Hughes in [22], [23] has generalized the concept of *group codes* for space-time codes. This class of codes is fully diverse and has distance invariance property. However, the rate of the code cannot achieve the maximum possible value (for full diversity, the maximum achievable rate is one symbol PCU [4].) The best known work in the category of non-orthogonal codes is [1] [8]. Here, an algebraic structure is utilized to construct non-orthogonal codes with a large coding advantage. For decoding, the authors suggest the *sphere decoding* algorithm (see [24]) that is generally a complex operation. Regarding the above statements, finding codes with good distance properties which at the same time allow the use of a simple Maximum Likelihood (ML) decoding algorithm remains an open problem.

In this paper, a new full-diversity full-rate space-time code is introduced. In this method, each code is equal to the addition of two matrices; \mathbf{A}^m and $\mathbf{D}\mathbf{A}^n$, where m and n are two data symbols and \mathbf{A} and \mathbf{D} are two constant matrices. \mathbf{A} is selected such that the set $\mathbf{A}^m, 0 \leq m \leq 2^b - 1$ is closed under matrix multiplication. This structure allows a simple Maximum Likelihood (ML) decoding method, and at the same time, simplifies the optimization of the coding advantage. This code is over two transmit antennas which has a significant practical importance and has been the subject of numerous papers [1], [15], [25]–[27]. Simulations show that the performance of the new code is very close to that of the Damen code [1] which is the best known block space-time code in terms of the coding advantage. Moreover, the decoding complexity of the proposed method is significantly less than that of the Damen code.

The rest of this paper is organized as follows: In Section II,

the system model and the basic design criteria are introduced. The structure of the proposed code is presented in Section III. In addition, a method to optimize the coding advantage is explained. In Section IV, a novel algorithm for the decoding of the proposed method is developed. In Section V, the simulation results and comparisons with the best alternative method are presented. Finally, the complexity of the proposed decoding method is compared with that of the Damen code which is based on sphere decoding.

II. PRELIMINARIES

In a slow flat fading environment, Multiple Input Multiple Output (MIMO) channels with M transmit antennas and N receive antennas over T symbol periods is modeled by

$$\mathbf{Y} = \sqrt{\frac{\rho}{M}} \mathbf{H} \mathbf{X} + \mathbf{V}, \quad (1)$$

where $\mathbf{Y} \in \mathcal{C}^{N \times T}$ denotes the received matrix, $\mathbf{X} \in \mathcal{C}^{M \times T}$ represents the transmitted matrix, $\mathbf{H} \in \mathcal{C}^{N \times M}$ signifies the channel matrix, and $\mathbf{V} \in \mathcal{C}^{N \times T}$ is the additive, spatially and temporally i.i.d white noise with a zero-mean unit-variance complex Gaussian distribution. Matrix \mathbf{X} , selected from a finite set, the *codebook*, and noise \mathbf{V} are normalized such that ρ is the SNR at each receive antenna. Considering the two matrices \mathbf{X}_η and \mathbf{X}_μ in the codebook, we define $\mathbf{B}_{\eta\mu}$ as $\mathbf{B}_{\eta\mu} = (\mathbf{X}_\eta - \mathbf{X}_\mu)(\mathbf{X}_\eta - \mathbf{X}_\mu)^*$, where $(\mathbf{X}_\eta - \mathbf{X}_\mu)^*$ is the conjugate transpose of the matrix $(\mathbf{X}_\eta - \mathbf{X}_\mu)$. The following are the criteria for space-time code design [4]:

- **Rank Criterion:** The diversity advantage is defined as the minimum rank of the matrices $\mathbf{B}_{\eta\mu}$ over all the pairs of codewords in the codebook. In order to achieve full diversity, the matrices $\mathbf{B}_{\eta\mu}$ have to be full rank for all pairs in the codebook.
- **Determinant Criterion:** The coding advantage is defined as the minimum of $[\prod_{l=1}^{\tau} \lambda_l(\mathbf{B}_{\eta\mu})]^{1/\tau}$, where $\lambda_l(\mathbf{B}_{\eta\mu})$, $l = 1 \dots \tau$ are the nonzero eigenvalues of $\mathbf{B}_{\eta\mu}$. For a fully diverse code, the coding advantage is equal to the minimum of $[\det(\mathbf{B}_{\eta\mu})]^{1/\tau}$.

III. CODE STRUCTURE

In this section, we introduce a new full-rate full-diversity block space-time code with a simple decoding method and high coding advantage.

The proposed structure is based on the PSK constellation with 2^b points. The codebook is composed of the following elements:

$$C = \{\mathbf{C}(m, n) \mid \mathbf{C}(m, n) = \mathbf{A}^m + \mathbf{D}\mathbf{A}^n, \text{ for } m, n = 0, 1, \dots, 2^b - 1\}, \quad (2)$$

where m and n denote two data symbols,

$$\mathbf{A} = \begin{bmatrix} \theta_1 & 0 \\ 0 & \theta_r \end{bmatrix} \quad \mathbf{D} = \begin{bmatrix} 0 & \theta_s \\ 1 & 0 \end{bmatrix}, \quad (3)$$

and

$$\theta_n = \exp(j \frac{2\pi}{2^b} n). \quad (4)$$

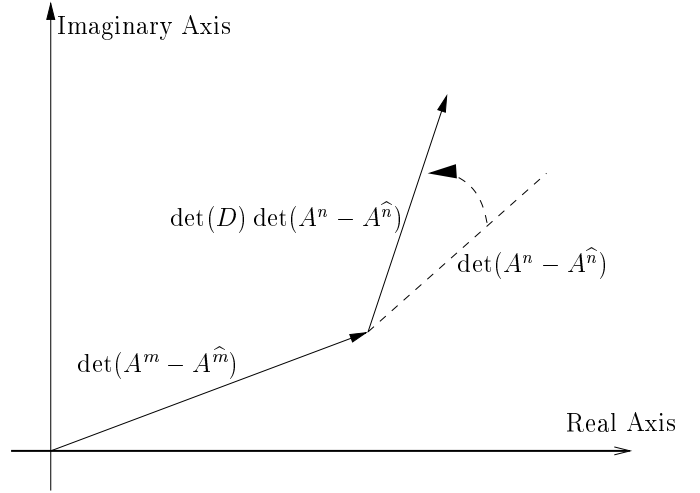


Fig. 1. Vector representation of $\det[\mathbf{C}(m, n) - \mathbf{C}(\hat{m}, \hat{n})]$.

Parameters r and s are selected for optimizing the coding advantage while maintaining the maximum diversity. The throughput of the code is equal to b bits or *one* symbol per channel use, the maximum achievable rate for a fully diverse code.

According to (2), each codeword $\mathbf{C}(m, n)$ is equal to the addition of two matrices namely \mathbf{A}^m and $\mathbf{D}\mathbf{A}^n$; each of them just depends on one of the data symbols m or n . In addition, it is evident that $\mathbf{A}^{-1} = \mathbf{A}^*$ and $\mathbf{A}^{2^b} = \mathbf{I}$, where \mathbf{I} is the identity matrix. In the next section, we show that this structure simplifies the Maximum Likelihood (ML) decoding. It also allows the use of a simple procedure for optimizing the coding advantage, as explained in the following.

Consider the following codebook:

$$C = \{\mathbf{C}(m, n) \in \mathcal{C}^{2 \times 2} \mid m, n = 0, 1, \dots, 2^b - 1\}. \quad (5)$$

The main objective is to maximize the coding advantage which is defined by (6) (shown at the top of the next page) for all the different pairs of (m, n) and (\hat{m}, \hat{n}) .

Using the basic matrix theory, we can rewrite the coding advantage as follows:

$$\Delta = \min_{m, n, \hat{m}, \hat{n}} |\det[\mathbf{C}(m, n) - \mathbf{C}(\hat{m}, \hat{n})]|, \quad (7)$$

where $|\cdot|$ denotes the magnitude of a complex number.

According to (2),

$$\det[\mathbf{C}(m, n) - \mathbf{C}(\hat{m}, \hat{n})] = \det(\mathbf{A}^m + \mathbf{D}\mathbf{A}^n - \mathbf{A}^{\hat{m}} - \mathbf{D}\mathbf{A}^{\hat{n}}). \quad (8)$$

Since $\mathbf{C}(m, n) - \mathbf{C}(\hat{m}, \hat{n})$ is a 2×2 matrix and since $\mathbf{A}^m - \mathbf{A}^{\hat{m}}$ is diagonal and also the main diagonal entries of $\mathbf{D}(\mathbf{A}^m - \mathbf{A}^{\hat{n}})$ are zero, it is obvious that

$$\det[\mathbf{C}(m, n) - \mathbf{C}(\hat{m}, \hat{n})] = \det(\mathbf{A}^m - \mathbf{A}^{\hat{m}}) + \det(\mathbf{D}) \det(\mathbf{A}^n - \mathbf{A}^{\hat{n}}). \quad (9)$$

By using the vector interpretation of the complex numbers, (9) is represented on the complex plane in Fig. 1. As it is shown in this figure, the vector $\det[\mathbf{C}(m, n) - \mathbf{C}(\hat{m}, \hat{n})]$ is equal to the addition of the two vectors: $\det(\mathbf{A}^m - \mathbf{A}^{\hat{m}})$ and $\det(\mathbf{D}) \det(\mathbf{A}^n - \mathbf{A}^{\hat{n}})$. Note that $\det(\mathbf{D}) = -\exp(js \frac{2\pi}{2^b})$.

$$\Delta = \min_{m,n,\hat{m},\hat{n}} \{ \det([\mathbf{C}(m,n) - \mathbf{C}(\hat{m},\hat{n})][\mathbf{C}(m,n) - \mathbf{C}(\hat{m},\hat{n})]^*) \}^{1/2} \quad (6)$$

TABLE I
OPTIMUM CODE STRUCTURE FOR 2^b PSK

	r	s	Coding Advantage
4PSK	1, 3	1, 3	2
8PSK	3	1, 3, 5, 7	1.0824
16PSK	7	2, 6, 10, 14	0.4483
32PSK	7, 23	3, 5, 11, 19, 21, 27, 29	0.1175

Therefore, $\det(\mathbf{D})$ is a vector with a magnitude equal to one and is used to change the direction of the second vector to increase the minimum length of $\det[\mathbf{C}(m,n) - \mathbf{C}(\hat{m},\hat{n})]$ which is the coding advantage.

Consequently, we can decouple the optimization procedure for selecting the parameters r and s , resulting in the following low complexity, suboptimum decoding procedure.

- *Step One:* The parameter r must be chosen to maximize the minimum of the $|\det(\mathbf{A}^m - \mathbf{A}^{\hat{m}})|$ for all different values of m and \hat{m} . (or $|\det(\mathbf{A}^n - \mathbf{A}^{\hat{n}})|$ for all the different values of n and \hat{n} .)

Since \mathbf{A} is a unitary matrix, $|\det(\mathbf{A}^m - \mathbf{A}^{\hat{m}})| = |\det(\mathbf{I} - \mathbf{A}^{(\hat{m}-m)})|$, indicating that the codebook $\{\mathbf{A}^m, m = 0, 1, \dots, 2^b - 1\}$ has distance invariance property. Therefore, optimizing the coding advantage is simplified to finding r to maximize the following function:

$$\min_{0 \leq m \leq 2^b - 1} |\mathbf{I} - \mathbf{A}^m| \quad (10)$$

which is the coding advantage of the codebook $\{\mathbf{A}^m, m = 0, 1, \dots, 2^b - 1\}$. This optimization has been accomplished and the results are reported in [22]. It is shown that the sets $\{\mathbf{A}^m, m = 0, 1, \dots, 2^b - 1\}$ and $\{\mathbf{D}\mathbf{A}^n, n = 0, 1, \dots, 2^b - 1\}$ are the *generalized Slepian's group code* [28], [22], [23].

- *Step Two:* In this step, we find the best value of s (or matrix \mathbf{D}) to maximize the coding advantage by using the selected value for r in step one. The matrix \mathbf{D} only changes the direction of the second vector in order to maximize the coding advantage. This rule of matrix \mathbf{D} is mathematically explained in Appendix I. However, the best value of s can be easily determined by a computer search.

In Table I, the best values for s and r for 2^b PSK, $2 \leq b \leq 5$ that optimize the coding advantage are listed.

IV. DECODING

To formulate the ML decoding, we have

$$P(\mathbf{Y}|\mathbf{H}, \mathbf{C}(m,n)) = \frac{1}{\pi^{NT}} \exp(-d^2(m,n)), \quad (11)$$

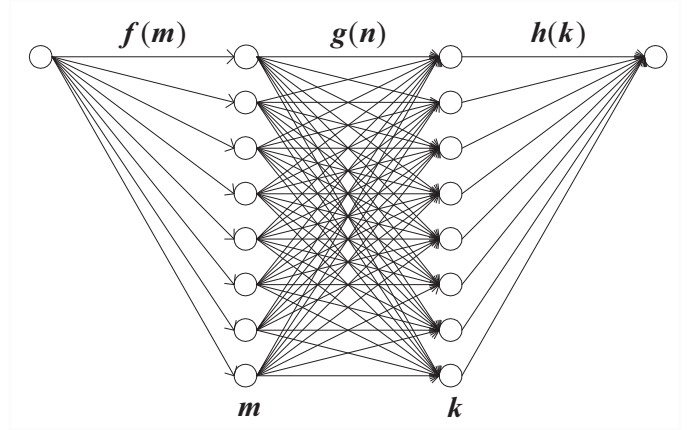


Fig. 2. Trellis construction based on code structure.

where \mathbf{H} is the channel transfer matrix, \mathbf{Y} is the matrix corresponding to the received signal, and

$$d^2(m,n) = \text{tr} \left[(\mathbf{Y} - \sqrt{\frac{\rho}{M}} \mathbf{H} \mathbf{C}(m,n)) (\mathbf{Y} - \sqrt{\frac{\rho}{M}} \mathbf{H} \mathbf{C}(m,n))^* \right], \quad (12)$$

and “tr” denotes trace function.

The goal of the ML decoding is to find m and n to maximize $P(\mathbf{Y}|\mathbf{H}, \mathbf{C}(m,n))$ or to minimize $d^2(m,n)$. The straightforward approach for the ML decoding is to calculate the different values of $d^2(m,n)$ for all the possible values of m and n , and to find the minimum value of $d^2(m,n)$ by using an exhaustive search. It is clear that such an exhaustive search is very complex.

Let us define K , $f(n)$, $g(m)$ and $h(m-n)$ as

$$K = \text{tr} \left\{ \mathbf{Y} \mathbf{Y}^* + 2 \frac{\rho}{M} \mathbf{H} \mathbf{H}^* \right\}, \quad (13)$$

$$f(m) = -\sqrt{\frac{\rho}{M}} \text{tr} \{ \mathbf{H} \mathbf{A}^m \mathbf{Y}^* + \mathbf{Y} \mathbf{A}^{-m} \mathbf{H}^* \}, \quad (14)$$

$$g(n) = -\sqrt{\frac{\rho}{M}} \text{tr} \{ \mathbf{H} \mathbf{D} \mathbf{A}^n \mathbf{Y}^* + \mathbf{Y} \mathbf{A}^{-n} \mathbf{D}^* \mathbf{H}^* \}, \quad (15)$$

and

$$h(m-n) = \frac{\rho}{M} \text{tr} \{ \mathbf{H} \mathbf{A}^{m-n} \mathbf{D}^* \mathbf{H}^* + \mathbf{H} \mathbf{D} \mathbf{A}^{n-m} \mathbf{H}^* \}. \quad (16)$$

By using these notations, it is easy to show that

$$d^2(m,n) = K + f(m) + g(n) + h(m-n). \quad (17)$$

To prove (17), we use these facts that

$$\mathbf{A}^* = \mathbf{A}^{-1} \quad \text{and} \quad \mathbf{D}^* = \mathbf{D}^{-1}. \quad (18)$$

Since $\mathbf{A}^{2^b} = \mathbf{A}^0 = \mathbf{I}_{2 \times 2}$, $h(m-n) = h(m-n \bmod 2^b)$. Thus, there are 2^b different values for each of $f(m)$, $g(n)$, and $h(m-n)$, where $0 \leq m \leq 2^b - 1$ and $0 \leq n \leq 2^b - 1$.

Ignoring the constant part K , for minimizing $f(m) + g(n) + h(m-n)$, we can use the Viterbi algorithm over the trellis

TABLE II

CODING ADVANTAGE OF THE ALAMOUTI, DAMEN, AND PROPOSED CODES
IN 3 BITS PCU

Structure	Alamouti Code	Damen Code	Proposed Code
Coding Advantage	0.5857	0.8074	1.0824

TABLE III

CODING ADVANTAGE OF THE ALAMOUTI, DAMEN, AND PROPOSED CODES
IN 4 BITS PCU

Structure	Alamouti Code	Damen Code	Proposed Code
Coding Advantage	0.1522	0.4738	0.4483

structure which is exemplified in Fig. 2 for the case of 8-PSK. In this figure, $k = [(m - n) \bmod 2^b]$.

Since this trellis is fully connected, it is possible to simplify the Viterbi algorithm to another effective ML decoding method with less complexity as explained by the following.

Let us sort $f(m)$ and $g(n)$ in the increasing order and specify the corresponding arguments as $m_0, m_1 \dots m_{2^b-1}$ and $n_0, n_1 \dots n_{2^b-1}$, i.e.,

$$f(m_0) \leq f(m_1) \leq \dots \leq f(m_{2^b-1}) \quad (19)$$

and

$$g(n_0) \leq g(n_1) \leq \dots \leq g(n_{2^b-1}). \quad (20)$$

We define two sets, the *potential set* and the *final set*. Each set has 2^b entries corresponding to the different values of $[(m - n) \bmod 2^b]$. The k^{th} entries of the *final set* $0 \leq k \leq 2^b - 1$ is the best pair of (m, n) in terms of the minimization of $f(m) + g(n)$ such that $k = [(m - n) \bmod 2^b]$. The k^{th} entry of the *potential set* $0 \leq k \leq 2^b - 1$ is the best pair of (m, n) until then in terms of the minimization of $f(m) + g(n)$ such that $k = [(m - n) \bmod 2^b]$, and for those values of k that are not yet in the *final set*.

All 2^b rows in *final set* and *potential set* are empty at the beginning and are gradually filled by employing the following algorithm.

For $0 \leq \alpha \leq 2(2^b - 1)$, starting from $\alpha = 0$,

- 1) Find $S_\alpha = \{(p, q) | p + q = \alpha, 0 \leq p \leq 2^b - 1, 0 \leq q \leq 2^b - 1\}$, where p and q are two integer numbers.
- 2) Find $(p, q) \in S_\alpha$ that minimize $f(m_p) + g(n_q)$.
- 3) Set $k = [(m_p - n_q) \bmod 2^b]$. Compare (m_p, n_q) with the pair in the k^{th} row of the *potential set*, if any, in terms of $f(m) + g(n)$, and insert the pair with the smaller value of $f(m) + g(n)$ in the k^{th} row of the *final set* (if the k^{th} row of the *final set* is not filled yet).
- 4) For the other pairs of $(p, q) \in S_\alpha$, compute $k = [(m_p - n_q) \bmod 2^b]$ and compare $f(m_p) + g(n_q)$ with the related value in the k^{th} row of the *potential set*. Substitute the pair with the smaller value of $f(m) + g(n)$ in the k^{th} row of the *potential set*.
- 5) If the *final set* is not filled yet and $\alpha < 2(2^b - 1)$, set $\alpha \leftarrow \alpha + 1$ and return to the first step, or else fill the empty rows of the *final set* with the *potential set*.
- 6) Compute $f(m) + g(n) + h(m - n)$ for the pairs in the *final set* and select the pair that minimizes $f(m) + g(n) + h(m - n)$.

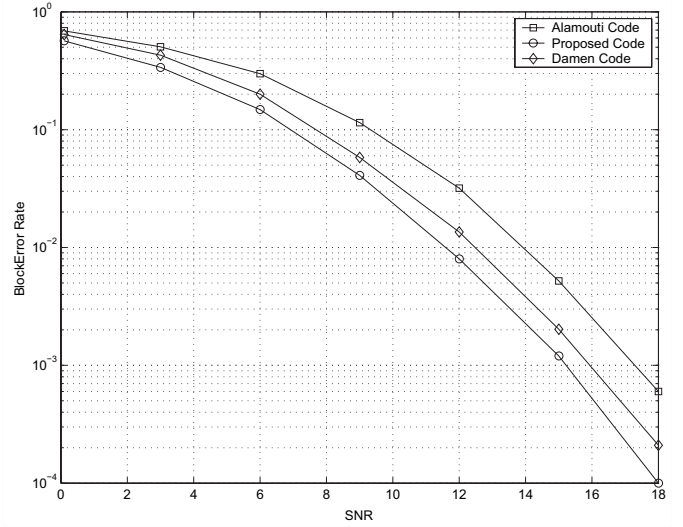


Fig. 3. Block error rate for two transmit and two receive antennas, 3 bits PCU.

The simulations demonstrate that the truncation of the algorithm in $\alpha = 2^b - 1$ does not considerably affect the code performance.

The idea behind the algorithm is as follows. If $(p, q) \in S_\alpha$ minimize $f(m_p) + g(n_q)$ and $k = [(m_p - n_q) \bmod 2^b]$, then there is no $(p', q') \in S_{\alpha'}$, $(\alpha < \alpha')$ such that $k = [(m_{p'} - n_{q'}) \bmod 2^b]$ and $f(m_p) + g(n_q) > f(m_{p'}) + g(n_{q'})$. The complete proof of the algorithm is in Appendix II.

With some modifications, we can use this method to decode BLAST systems (see [3]) with two transmit antennas [29].

V. COMPARISON AND SIMULATION RESULTS

In this section, an evaluation of the code performance by means of simulations to measure the block error rate is presented. In addition, we compare the performance of this code with that of two well-known codes, so called Alamouti code [15] and Damen code [1]. This comparison is presented for the case of three and four bit PCU. In the following comparisons, we use the optimal ML method for decoding; however, in the next section, we compare the performance of these codes considering the effects of suboptimal decoding.

In the case of a three bits PCU, we use 8-PSK modulation. Table II indicates the coding advantages of the Alamouti code, the Damen Code and the proposed code for a three bits throughput. It is clear that there is significant improvement of the coding advantage of the proposed code compared to the coding advantage of the Alamouti code. In addition, the performance of the proposed code is close to the performance of the Damen code.

We now look at the block error rate curves of these two codes. Fig. 3 denotes the block error rate curves for the Alamouti code, the Damen code, and proposed scheme. This figure confirms that there is more than 2dB improvement in the coding gain as compared to the Alamouti code, and about 1dB as compared to the Damen code.

For the 4 bits PCU, the 16PSK constellation is used for the modulation in the proposed code. Table III displays the coding advantage of the Damen scheme and the proposed scheme. It

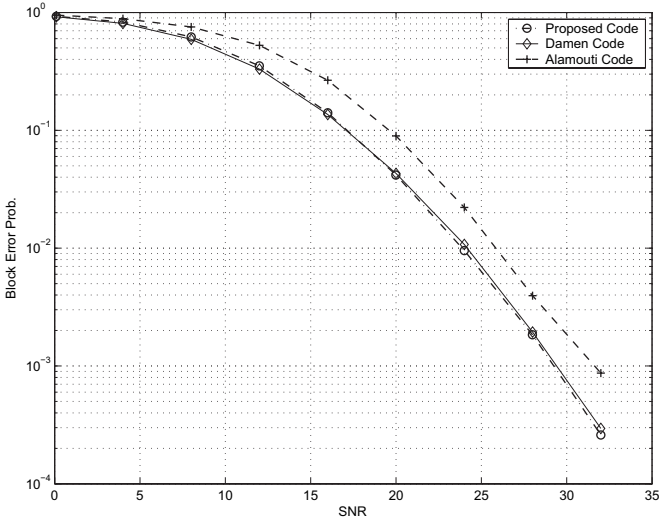


Fig. 4. Block error rate for two transmit and one receive antennas, 4 bits PCU.

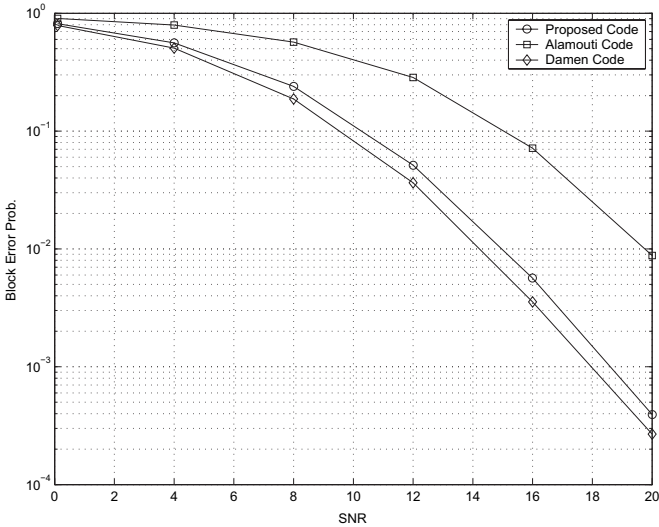


Fig. 5. Block error rate for two transmit and two receive antennas, 4 bits PCU.

is apparent that the coding advantage of the novel code is very close to the coding advantage of the Damen code.

Figures 4 and 5 depict the block error rate curves for one and two receive antennas, respectively. It is apparent that the performance of the proposed code is similar to that of the Damen code. In the next section, we will compare the complexity of the decoding of the proposed scheme with that of sphere decoding that is used for the decoding of the Damen code. In addition, we will compare the performances under some constraints regarding the complexity of the decoding.

The proposed code is based on PSK constellation. In PSK constellation with constant energy, the minimum distance of the constellation points decreases fast with increasing the size of constellation. Therefore, the PSK constellation is not recommended for very high spectral efficiency. As a results, the proposed code is appropriate for the rates in the range of less than or equal 4 bits PCU which is the case in most practical wireless systems.

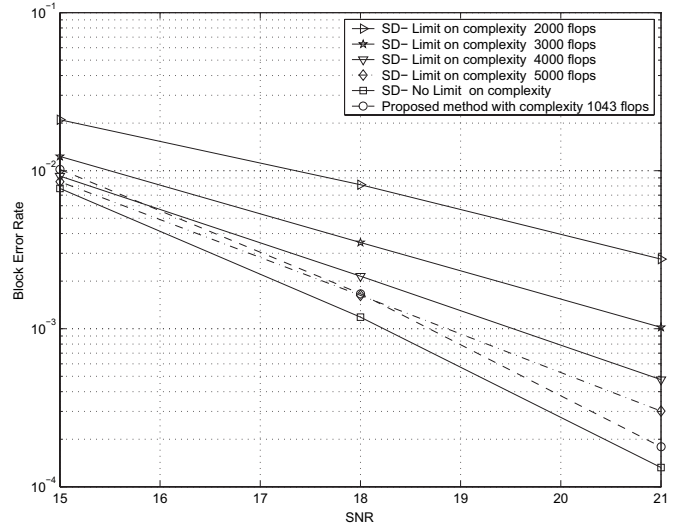


Fig. 6. Effect of the complexity truncation over Sphere Decoding (SD), as compared with the proposed method for two transmit and two receive antennas, 4 bits PCU.

On the other hand, using PSK constellation in the proposed structure results in a Peak-to-Average Ratio (PAR) [30] equal to one which is almost half of the PAR for the Damen code.

VI. COMPLEXITY

In this section, the complexity of the proposed decoding method is compared to that of the sphere decoding which is used for the Damen scheme. The maximum complexity of the proposed decoding method with a 2^b PSK constellation and N receiver antennas is equal to $2 \times 2^{2b} + (29 + 2b)2^b + 48N - 19$ flops, where the complexity is defined as the total numbers of additions, multiplications, and comparisons. For instance, the maximum complexity of the decoding method for 16PSK and two receive antennas is equal to 1181 flops. Although this equation shows the maximum complexity of the proposed decoding method, the simulations exhibit that the complexity for different realizations of channel matrix \mathbf{H} is almost constant.

On the other hand, the complexity of sphere decoding largely varies, depending on the channel matrix \mathbf{H} . As a result, in order to practically implement the sphere decoding, we need to truncate the decoding algorithm whenever the complexity is greater than a pre-determined threshold. This truncation has an undesirable effect on the performance of the code. Figure 6 shows the effect of the truncation on the performance of the Damen code in the 4 bits PCU. The implementation is based on the best known sphere decoding algorithm presented in [31] (*Algorithm II*). In [31], the Schonor-Euchner algorithm (see [32]) is modified to search the points inside a sphere centered at the received vector.

The radius of the sphere in which the search is accomplished is very important. If the radius is large, the sphere decoding will be very complex. On the other hand, if the radius is small, the additional error caused by decoding procedure will be very high. As a trade-off between complexity and additional error, we select the radius of the sphere such that the additional error is ten percent of the current error rate.

Figure 6 also shows that the performance of the proposed method is significantly improved as compared to the Damen code, if we truncate the algorithm so that the complexity is less than a given threshold. According to the curves in Fig. 6, it is easy to see that for a similar performance in the range of 10^{-2} to 10^{-3} block error rate, we have to truncate the decoding complexity of sphere decoding to at least 4000 flops.

Thus, when the performance is equal, the complexity of the proposed method is about four times less than the complexity of the Damen code (decoded using sphere decoding). Note that the limitation over the complexity of decoding sacrifices the diversity of the Damen code.

VII. CONCLUSION

In this paper, a new non-orthogonal full-diversity full-rate block space-time code over two transmit antennas and two symbol periods is introduced. In this method, each code is equal to the addition of two matrices; \mathbf{A}^m and $\mathbf{D}\mathbf{A}^n$, where m and n are the two data symbols and \mathbf{A} and \mathbf{D} are the two constant matrices. \mathbf{A} is selected such that the set $\mathbf{A}^m, 0 \leq m \leq 2^b - 1$ is closed under matrix multiplication. This structure allows a simple Maximum Likelihood (ML) decoding method, and at the same time, simplifies the optimization of the coding advantage. A new method of decoding with a low complexity is also developed, showing about four times reduction in the decoding complexity (for similar performance) as compared to the best known alternative code introduced in [1].

APPENDIX I

AN EXPLANATION FOR DETERMINING THE PARAMETER s

As we mentioned in the paper, $\det[\mathbf{C}(m, n) - \mathbf{C}(\hat{m}, \hat{n})]$ has two parts: $\det(\mathbf{A}^m - \mathbf{A}^{\hat{m}})$ and $\det(\mathbf{D})\det(\mathbf{A}^n - \mathbf{A}^{\hat{n}})$. The role of the matrix \mathbf{D} is to rotate the second part such that the minimum norm of the addition of this two parts is maximum. In the following, the selection of the matrix \mathbf{D} (or parameter s) will be detailed. Let us start with the following equation:

$$\begin{aligned} & \det[\mathbf{C}(m, n) - \mathbf{C}(\hat{m}, \hat{n})] \\ &= \det(\mathbf{A}^m - \mathbf{A}^{\hat{m}}) + \det(\mathbf{D})\det(\mathbf{A}^n - \mathbf{A}^{\hat{n}}) \\ &= \left(e^{jm\frac{2\pi}{2^b}} - e^{j\hat{m}\frac{2\pi}{2^b}}\right) \left(e^{jn\frac{2\pi}{2^b}} - e^{j\hat{n}\frac{2\pi}{2^b}}\right) \\ &\quad - e^{js\frac{2\pi}{2^b}} \left(e^{jn\frac{2\pi}{2^b}} - e^{j\hat{n}\frac{2\pi}{2^b}}\right) \left(e^{jr\frac{2\pi}{2^b}} - e^{j\hat{r}\frac{2\pi}{2^b}}\right) \\ &= -4 \exp \left[j(m + \hat{m})(r + 1)\frac{\pi}{2^b} \right] \\ &\quad \times \sin \left[\frac{\pi}{2^b}(m - \hat{m}) \right] \sin \left[r\frac{\pi}{2^b}(m - \hat{m}) \right] \\ &\quad - 4 \exp \left(js\frac{2\pi}{2^b} \right) \exp \left[j(n + \hat{n})(r + 1)\frac{\pi}{2^b} \right] \\ &\quad \times \sin \left[\frac{\pi}{2^b}(n - \hat{n}) \right] \sin \left[r\frac{\pi}{2^b}(n - \hat{n}) \right]. \end{aligned} \quad (21)$$

Equation (21) helps us to determine the best values for s . In the following, the procedure for determining s for 8PSK, as an example, will be explained. However, it is easy to generalize this method to a PSK constellation with a different number of points. According to (21), for the 8PSK constellation ($b = 3$),

the following can be written:

$$\begin{aligned} & \det[\mathbf{C}(m, n) - \mathbf{C}(\hat{m}, \hat{n})] \\ &= -4 \times j^{(m+\hat{m})} \times \sin \left[\frac{\pi}{8}(m - \hat{m}) \right] \sin \left[3\frac{\pi}{8}(m - \hat{m}) \right] \\ &\quad - 4 \times \exp \left(js\frac{2\pi}{8} \right) \times j^{(n+\hat{n})} \sin \left[\frac{\pi}{8}(n - \hat{n}) \right] \sin \left[3\frac{\pi}{8}(n - \hat{n}) \right]. \end{aligned} \quad (22)$$

Using the vector interpretation of complex numbers, $j^{(m+\hat{m})} \times \sin \left[\frac{\pi}{8}(m - \hat{m}) \right] \sin \left[3\frac{\pi}{8}(m - \hat{m}) \right]$ is a vector only on the directions of the real and imaginary axes of the complex plane. The same is true for $j^{(n+\hat{n})} \times \sin \left[\frac{\pi}{8}(n - \hat{n}) \right] \sin \left[3\frac{\pi}{8}(n - \hat{n}) \right]$. In this case, these two vectors can cancel each other and the coding advantage will be zero. To avoid this and to maximize the coding gain, the second vector needs to be rotated. The best angles for the rotation are 45, 135, 225, and 315 degrees. Therefore, s can be 1, 3, 5, and 7.

APPENDIX II

PROOF OF THE PROPOSED DECODING ALGORITHM

In this part, we prove the algorithm that is presented for decoding the proposed code. Consider S_α as defined in the first step of the algorithm. We define $(p_{min}(\alpha), q_{min}(\alpha))$ as follows:

$$(p_{min}(\alpha), q_{min}(\alpha)) = \arg \min_{(p,q) \in S_\alpha} \{f(m_p) + g(n_q)\}. \quad (23)$$

It is easy to show that if $\alpha' > \alpha$, then

$$\exists (p, q) \in S_\alpha \text{ such that } p \leq p_{min}(\alpha') \text{ and } q \leq q_{min}(\alpha'). \quad (24)$$

According to the inequalities (19) and (20), the inequality (24) results in

$$\begin{aligned} \exists (p, q) \in S_\alpha & : f(m_p) + g(n_q) \\ & \leq f(m_{p_{min}(\alpha')}) + g(n_{q_{min}(\alpha')}). \end{aligned} \quad (25)$$

On the other hand, $(p, q) \in S_\alpha$; therefore based on the definition of $(p_{min}(\alpha), q_{min}(\alpha))$,

$$f(m_{p_{min}(\alpha)}) + g(n_{q_{min}(\alpha)}) \leq f(m_p) + g(n_q). \quad (26)$$

By using (25) and (26), it easy to see that,

$$f(m_{p_{min}(\alpha)}) + g(n_{q_{min}(\alpha)}) \leq f(m_{p_{min}(\alpha')}) + g(n_{q_{min}(\alpha')}). \quad (27)$$

According to the definition of $(p_{min}(\alpha'), q_{min}(\alpha'))$ and considering (27), we can see that:

$$\begin{aligned} \forall (p, q) \in S_{\alpha'} & : f(m_{p_{min}(\alpha)}) + g(n_{q_{min}(\alpha)}) \\ & \leq f(m_p) + g(n_q). \end{aligned} \quad (28)$$

So, for any (p, q) in $S_{\alpha'}$ such that $(m_{p_{min}(\alpha)} - n_{q_{min}(\alpha)} \bmod 2^b) = (m_p - n_q \bmod 2^b)$,

$$\begin{aligned} f(m_{p_{min}(\alpha)}) &+ g(n_{q_{min}(\alpha)}) + h(m_{p_{min}(\alpha)} - n_{q_{min}(\alpha)}) \\ &\leq f(m_p) + g(n_q) + h(m_p - n_q). \end{aligned} \quad (29)$$

Inequality (29) guarantees that if we define $k = (m_{p_{min}(\alpha)} - n_{q_{min}(\alpha)} \bmod 2^b)$, then the k^{th} row of the final set should be filled with $(m_{p_{min}(\alpha)}, n_{q_{min}(\alpha)})$, and there is no better choice for (p, q) in $S_{\alpha'}$ ($\alpha' > \alpha$).

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